

Stochastic Approach to CSP's in the Hardware Verification Domain

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Simulation Based Methods, Systems and Modeling

Constraints in test case generation

Verification task:
Real address in some corner memory space,

Effective address aligned to 64K.

Testing knowledge: Reuse cache row.

Architecture: Effective address translates into real address in a complex

way.

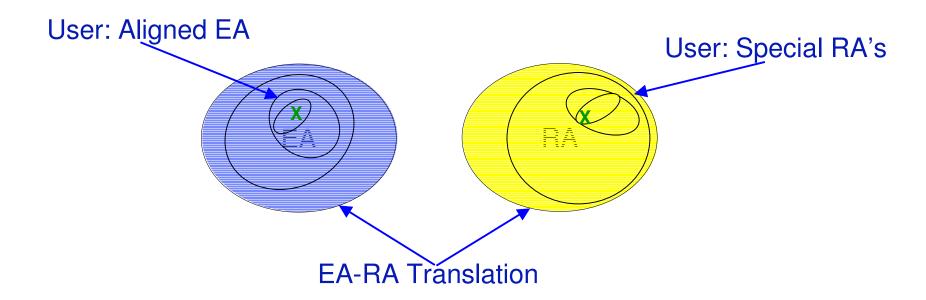
EA: $0 \times 0 B = 274 \text{ FAB} = 0 D B \in 00000$ Huge Domains (2⁶⁴)

RA: 0x0002FFC5_90A4D000 Quality of solution: Random uniformity

EA: 0x0002FF00_0000000

RA: 0x0002FF00_00000000 Correct, but worthless!

Deterministic Methods (DPLL, MAC, k-consistency, ...)



- 1. Reach a level of consistency through successive reductions of sets
- 2. Choose a random assignment for a variable, and maintain the consistency

Limitations of Deterministic Methods: An example

$$a,b,c \in \{0,...,N\}, N = 2^{64}$$

1.
$$a \neq 0 \Rightarrow b = 0$$

2. $a = 0 \Rightarrow c = 0$
3. $b = 0 \Rightarrow c = 0$
4. $c = 0 \Rightarrow a = 1$
Only solution: $a = 1, b = c = 0$

Local consistency at onset: Choose randomly with probability 1/N of being correct (Solution reached at 600 million years)

Limitations of Deterministic Methods: Another example

$$a,b,c \in \{0,...,N\}, N = 2^{64}$$

1.
$$a*b=c$$

2. a,b,c each have five 1's in their binary representation

Already a single reduction of domains is hard

Stochastic approaches: defining the metrics

- State: a tuple representing a single assignment to each variable
- Cost: A function from the set of states to {0} U R⁺
 - \square Cost = 0 iff all constraints are satisfied by the state.

$$a*b=c$$
 $\operatorname{Cost} = (c \square a*b)^2$

$$a = 0 \Rightarrow b = 0$$
 Cost $= \begin{bmatrix} b & a = 0 \\ 0 & a \neq 0 \end{bmatrix}$

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1.2